# **Prolog Basics**

### CS3100 Fall 2019

### **Review**

### Previously

• Functional Programming in OCaml

### Today

Introduction to Logic Programming in Prolog

### Imperative programming

Computing the sum of the elements of an integer list in Java.

```
int sum (int[] list) {
    int result = 0;
    for (int i = 0; i < list.length; i++)
        result += list[i];
    return result;
}</pre>
```

## **Functional Programming**

Computing the sum of the elements of an integer list in OCaml.

let rec sum 1 = match 1 with
 [] -> 0
 [ x::xs -> x + sum xs

## **Logic Programming**

Computing the sum of the elements of an integer list in **Prolog**.

In [1]:

sum([],0).
sum([H | T], N) :- sum(T,M), N is H+M.

```
Added 2 rule(s).
```

Notice that this is a **declarative** reading of the sum of a list.

## **Declarative vs Operational**

- This Prolog program says what the sum of a list is.
  - OCaml and Java programs were about **how** to compute the sum.
- In particular, prolog program does not define **control flow** through the program.
  - program is a collection of facts and rules

### **Prolog Program Answers Questions**

+----+
Queries ==> | Facts + Rules | ==> Answers
+----+
Prolog Program

Facts and rules together build up a database of relations.

### Relational view of the sum program

The program

sum([],0)
sum([H | T], N) :- sum(T,M), N is M+H

inductively defines a table of relations:

++			
L	ist	Sum	
[	]	0	
[	1]	1	
[	1,2]	3	
[	2]	2	
.	••	• • •	

# Queries are look ups in this table

### In [2]:

```
?- sum([1,2,3],X).
```

X = 6.

Of course, the **computation model** is not to build a database and look up facts.

## Why this declarative view?

- Many problems in computer science are naturally expressed as declarative programs.
  - Rule-based AI, Program Analysis (asking questions on code), Type Inference, queries on graphical programs, UIs.
- But the programmer has to convert this to Von Neumann Architecture.



# Logic Programming to the rescue

- Logic programming the programmer to declaratively express the program
- The compiler will figure out how to compute the answers to the queries.

```
Prolog = Logic (programmer) + Control (compiler)
```

# Prolog

- Is one of the first logic programming languagues
- to be precise, it is a family of languages that differ by the choice of control
- Invented in 1972, and has many different implementations
  - We will use **SWI-Prolog** for our study.

## **House Stark**



# **Prolog Terms**

Prolog programs are made up of terms.

- **Constants**: 1,2,3.14,robb,'House Stark', etc.
  - also known as atoms.
- Variables: Always begin with a capital letter.
  - X, Y, Sticks, \_.
- compound terms: male(robb), father(ned,robb).
  - Top-function symbol/functor: male, father
  - **arity**: Number of arguments; male = 1, father = 2.

• top function symbols also written down explicitly with arity such as male/1, father/2.

### House Stark -- Facts

#### In [3]:

```
father(rickard,ned).
father(rickard,brandon).
father(rickard,lyanna).
father(ned,robb).
father(ned,sansa).
father(ned,arya).
```

Added 6 rule(s).

### **House Start -- Queries**

#### In [4]:

```
?- father(ned,sansa).
```

true.

#### In [5]:

?- father(rickard,sansa).

false.

### **Closed world assumption**

We know that Ned is the father of Bran.

Let us query our program.

#### In [6]:

```
?- father(ned,bran).
```

false.

• Closed World Assumption: Prolog only knows the fact that it has been told.

- Assumes false for everything else.
- Interesting interactions with negation (we will see this later).

## **Existential Queries**

• Apart from true/false questions, we can also ask queries that return other answers (existential queries).

"Who are Ned's children?"

#### In [7]:

```
?- father(ned, X).
```

X = robb ; X = sansa ; X = arya .

## **Existential Queries**

"Who is the father of Arya?"

#### In [8]:

```
?- father(X,arya).
```

X = ned.

"Who are Robb's children?"

```
In [9]:
```

?- father(robb,X).

false.

### **Rules**

- So far what we have done could have been done with a relational database.
- Rules define further facts inductively from other facts and rules.
- Rules have a head and body.
  - H :- B1, B2, B3, ..., BN

• H is true if  $B1 \wedge B2 \wedge B3 \dots BN$  is true.

## **Rules**

#### In [10]:

```
parent(X,Y) := father(X,Y).
```

```
ancestor(X,Y) :- parent(X,Y).
ancestor(X,Y) :- parent(X,Z), ancestor(Z,Y).
```

```
Added 3 rule(s).
```

Observe that z only appears on the RHS of the last rule.

# Rules

#### In [11]:

```
?- ancestor(rickard,X).
```

X = ned ; X = brandon ; X = lyanna ; X = robb ; X = sansa ; X = arya .

### **Exercise**

Define mother, cousin, uncle, aunt, sibling.

## Quiz

#### In [12]:

```
material(gold).
material(aluminium).
process(bauxite,alumina).
process(alumina,aluminium).
process(copper, bronze).
valuable(X) :- material(X).
valuable(X) :- process(X,Y), valuable(Y).
```

```
Added 7 rule(s).
```

- Which of these are valuable?
  - gold, bauxite, bronze, copper.

### Quiz

#### In [13]:

```
?- valuable(gold).
?- valuable(bauxite).
?- valuable(bronze).
?- valuable(copper).
```

true. true. false. false.

## Unification

At the core of how Prolog computes is Unification.

There are 3 rules for unification:

- · Atoms unify if they are identical
  - a & a unify, but not a & b.
- Variables unify with anything.
- Compound terms unfig only if their top-function symbols and arities match and their arguments unify recursively.

Quiz

Which of these unify?

- 1. a & a
- 2. a & b
- 3. a & A
- 4. a & B
- 5. tree(l,r) & A

## Quiz

Which of these unify?

1. a & a yes
 2. a & b no
 3. a & A yes
 4. a & B yes
 5. tree(l,r) & A yes

# Quiz

Which of these unify?

- 1. tree(I,r) & tree(B,C)
- 2. tree(A,r) & tree(I,C)
- 3. tree(A,r) & tree(A,B)
- 4. A & a(A)
- 5. a & a(A)

# Quiz

Which of these unify?

- 1. tree(I,r) & tree(B,C) yes
- 2. tree(A,r) & tree(I,C) yes
- 3. tree(A,r) & tree(A,B) **yes**
- $4.\ A \ \& \ a(A)$  yes (mostly), occurs check disabled by default
- 5. a & a(A) no

# Note about prolog notebooks

- There are no binders for the rules and facts in prolog.
  - This is unlike OCaml which has top-level definitions.
- Hence, you may see strange behaviours when working with the notebook.

# Note about notebooks

### In [14]:

```
stringofint(1,"one").
```

Added 1 rule(s).

In [15]:

?- stringofint(1,X).

X = one.

Restart the kernel and run again if you find weird results.

# Fin.